Lecture 13: Transport Layer Flow and Congestion Control COMP 332, Fall 2018 Victoria Manfredi





Acknowledgements: materials adapted from Computer Networking: A Top Down Approach 7th edition: ©1996-2016, J.F Kurose and K.W. Ross, All Rights Reserved as well as from slides by Abraham Matta at Boston University, and some material from Computer Networks by Tannenbaum and Wetherall.

Today

1. Announcements

– exam wed!

2. Midterm overview

exam format

3. TCP

- reliable data transfer
- connection management
- flow control

Midterm OVERVIEW

Midterm overview

In class on Wednesday, Oct. 17

- closed book, closed notes
- covers material in lectures 1 to 12

Will not ask questions on

- probability

5 questions

- app layer short questions
- HTTP persistent vs. non-persistent connections
- transport layer short questions
- socket coding
- reliable data transport protocol

Problems 1 and 3

App layer and transport layer short questions

- 8 in total
- similar to review questions in book
- should only need to write a few sentences to answer

Problem 2

HTTP persistent vs. non-persistent connections

- review related homework question

Problems 4

Socket coding

- be able to write code to open, use, and close sockets
- differences between client and server code

Problem 5

Design a reliable data transfer protocol

- given channel characteristics design most efficient protocol
- be able to design reliable data transfer protocol like stop-and-wait
- know your timeline diagrams

TCP RELIABLE DATA TRANSFER

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Duplicate ACKs

Time-out period often relatively long

long delay before resending lost packet

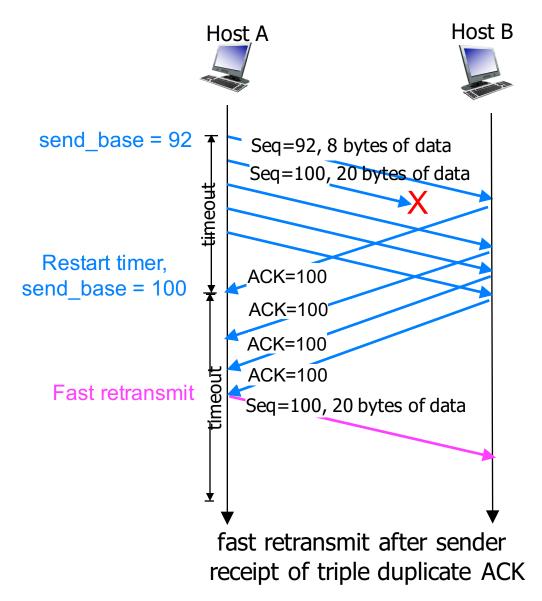
Duplicate ACKs indicate isolated loss

- rather than congestion causing many losses
 - sender often sends many segments back-to-back
 - if segment is lost, likely many duplicate ACKs
 - ACKs being received indicates some packets received at destination since ACK sent for every packet: so not congestion

TCP fast retransmit

- if sender receives 3 ACKs for same data (triple duplicate ACKs)
 - resend unacked segment with smallest seq #
- Q: why 3?
 - pkts may just have been reordered otherwise
 - likely that unacked segment lost, so don't wait for timeout

TCP fast retransmit



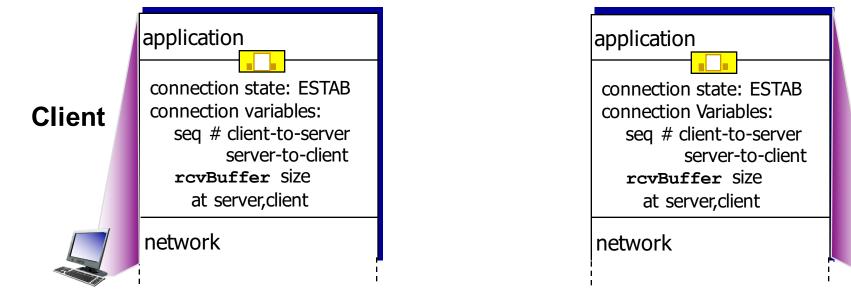
TCP CONNECTION MANAGEMENT

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Connection Management

Before exchanging data, sender/receiver handshake

- establish connection and connection parameters
- tear down connection when done



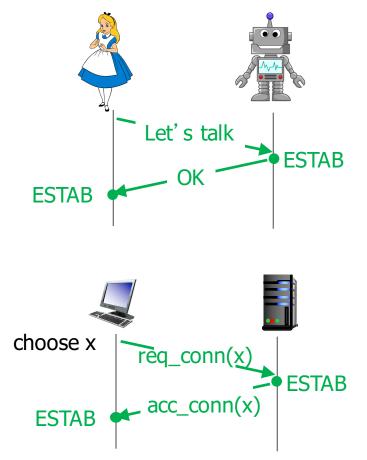
sock = sock.connect((host, port))

conn, addr = server_sock.accept()

Server

Agreeing to establish a connection

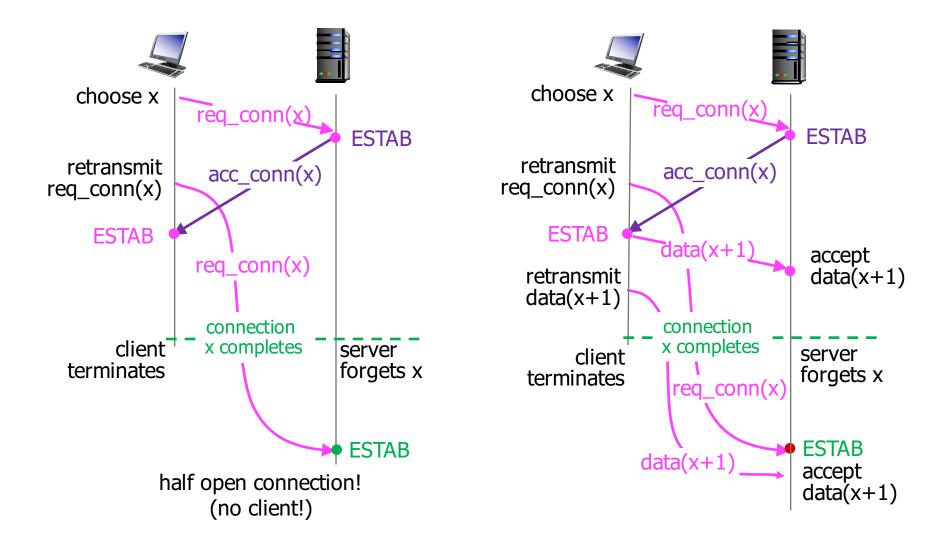
2-way handshake:



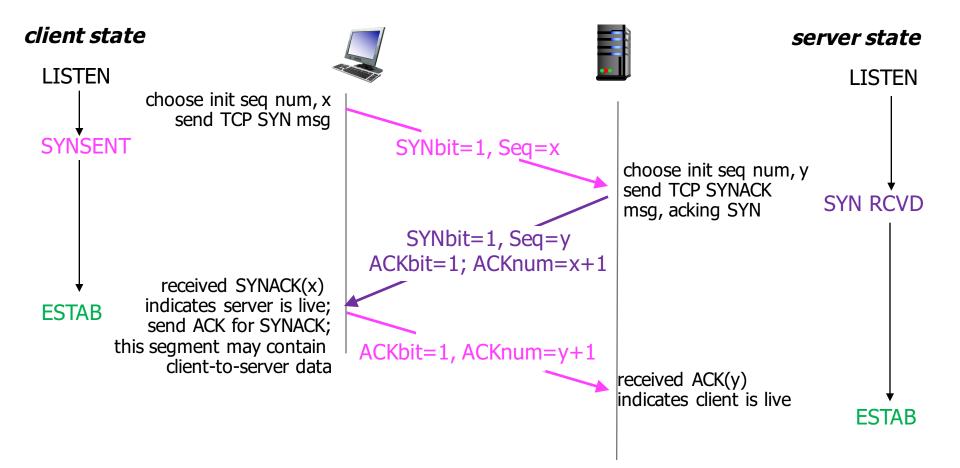
Q: will 2-way handshake always work in network?

- variable delays
- retransmitted messages
 - e.g. req_conn(x)) due to message loss
- message reordering
- can't see other side

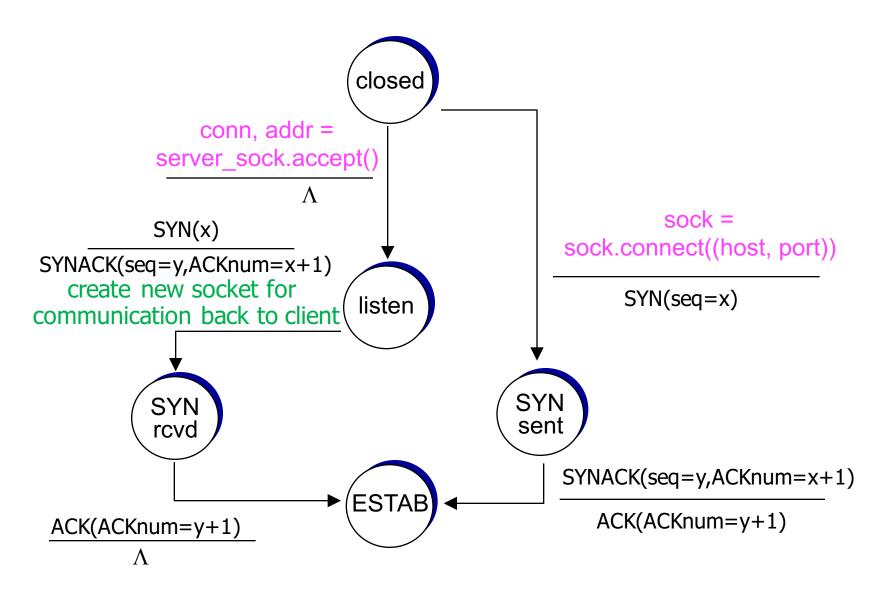
2-way handshake failure scenarios



TCP 3-way handshake



TCP 3-way handshake: FSM



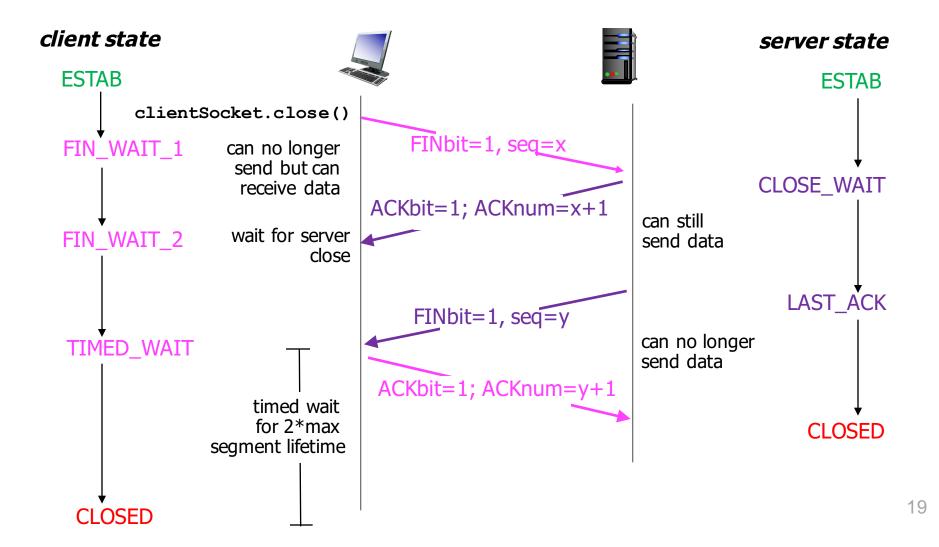
Look at the state of tcp connections

> netstat -ta							
Active	ctive Internet connections (including servers)						
Proto	Recv-Q Send	-Q	Local Address	Foreign Address	(state)		
tcp4	0	0	<pre>vmanfredismbp2.w.55777</pre>	lga25s60-in-f5.1.https	ESTABLISHED		
tcp4	31	0	<pre>vmanfredismbp2.w.55736</pre>	162.125.34.6.https	CLOSE_WAIT		
tcp4	0	0	<pre>vmanfredismbp2.w.55717</pre>	a104-110-151-148.https	ESTABLISHED		
tcp4	0	0	<pre>vmanfredismbp2.w.55716</pre>	a104-110-151-148.https	ESTABLISHED		
tcp4	0	0	<pre>vmanfredismbp2.w.55715</pre>	a104-110-151-148.https	ESTABLISHED		
tcp4	0	0	<pre>vmanfredismbp2.w.55714</pre>	a104-110-151-148.https	ESTABLISHED		
tcp4	0	0	<pre>vmanfredismbp2.w.55713</pre>	a104-110-151-148.https	ESTABLISHED		
tcp4	0	0	<pre>vmanfredismbp2.w.55668</pre>	wesfiles.wesleya.http	CLOSE_WAIT		
tcp4	0	0	<pre>vmanfredismbp2.w.55486</pre>	162.125.18.133.https	ESTABLISHED		
tcp4	0	0	<pre>vmanfredismbp2.w.55322</pre>	162.125.18.133.https	ESTABLISHED		
tcp4	31	0	<pre>vmanfredismbp2.w.55250</pre>	162.125.4.3.https	CLOSE_WAIT		
tcp4	0	0	<pre>vmanfredismbp2.w.55170</pre>	ec2-52-20-75-192.https	CLOSE_WAIT		
tcp4	0	0	<pre>vmanfredismbp2.w.55072</pre>	85.97.201.35.bchttps	ESTABLISHED		
tcp4	0	0	localhost.ipp	*.*	LISTEN		
tcp6	0	0	localhost.ipp	*.*	LISTEN		
tcp4	0	0	<pre>vmanfredismbp2.w.53453</pre>	6.97.a86c.ip4.st.https	ESTABLISHED		

TCP: politely closing a connection

Client, server each sends TCP segment with FIN bit = 1

- respond to received FIN with ACK (ACK can be combined with own FIN)



FIN segment in Wireshark

241 4.063493 vmanfredismbp2.wireless.we 40.97.120.226	54 55017 → 443 [FIN										
Frame 241: 54 bytes on wire (432 bits), 54 bytes captured (432 bits) on interface 0											
Ethernet II, Src: 78:4f:43:73:43:26 (78:4f:43:73:43:26), Dst: 129.133.176.1 (3c:8a:b0:1e:18:01)											
▶ Internet Protocol Version 4, Src: vmanfredismbp2.wireless.wesleyan.edu (129.133.187.174),	Dst: 40.97.120.226 (40.97.12										
Transmission Control Protocol, Src Port: 55017 (55017), Dst Port: 443 (443), Seq: 3771, Ac	ck: 6504, Len: 0										
Source Port: 55017											
Destination Port: 443											
[Stream index: 5]											
[TCP Segment Len: 0]											
Sequence number: 3771 (relative sequence number)											
Acknowledgment number: 6504 (relative ack number)											
<pre>Header Length: 20 bytes Flags: 0x011 (FIN, ACK) 000 = Reserved: Not set0 = Nonce: Not set0 = Congestion Window Reduced (CWR): Not set</pre>											
						\dots $0.\dots = ECN-Echo: Not set$					
						1 = Acknowledgment: Set					
$\dots \dots $											
0 = Reset: Not set											
$\dots \dots $											
1 = Fin: Set											
[TCP Flags: *****A***F]											
Window size value: 8192											
[Calculated window size: 262144]											
[Window size scaling factor: 32]											
Checksum: 0xe59d [validation disabled]											
0000 3c 8a b0 1e 18 01 78 4f 43 73 43 26 08 00 45 00 <x0 csc&e.<="" td=""><td></td></x0>											
0010 00 28 76 59 40 00 40 06 e5 ff 81 85 bb ae 28 61 .(vY@.@(a											
0020 78 e2 d6 e9 01 bb dd 11 e8 4a b0 93 7d 29 50 11 x											
0030 20 00 e5 9d 00 00											